

In preparing your solutions to the exam, you are **allowed to use any sources** including your textbook, other students and professors, previous homeworks and solutions, or any sources on the Internet. You may ask me for feedback on potential solutions, but I will not give you any hints. Of course, I am happy to answer general questions, go over homework problems, or answer clarifying questions about exam problems.

The exam will take place in class on Friday, February 24 (MC Reynolds 317, 8:10-9:00am). You are **not** allowed to bring any notes, textbooks, calculators, or any other resources with you to the exam. Bring only something to write with; I will provide a fresh copy of the exam, paper for writing your solutions, and scratch paper.

Question 1 (20 points). Characterize the asymptotic behavior of each of the following in terms of big-Theta, as a function of n and/or m . Prove/justify your answers. Full credit will only be given for the best (fastest) possible algorithms.

- (a) $1 + 2 + 3 + \dots + n/2$
- (b) $4 + 8 + 16 + 32 + \dots + 2^n$
- (c) Time to find the shortest path between two given vertices in an unweighted, undirected graph with n vertices and m edges.
- (d) Time to find the shortest path between two given vertices in a weighted, undirected graph (assuming all weights are nonnegative) with n vertices and m edges.
- (e) Given an array A of length n , the time to fill in a matrix M such that $M[i, j] = A[i] + A[j]$.
- (f) Given an array A containing n positive integers, the time to compute the smallest positive integer which is not contained in the array.
- (g) Given an array A containing n positive integers which are all at most 100, the time to compute the smallest positive integer which is not contained in the array.

Question 2 (15 points). Prove, or disprove with a counterexample: if a tree T has a vertex of degree $k \geq 1$, then T has at least k leaves.

Question 3 (15 points). Consider the following algorithm to determine whether an undirected, unweighted graph G has any cycles: pick an arbitrary vertex v and run a breadth-first search (BFS), generating a sequence of layers L_0, L_1, L_2, \dots . If there are any edges between two vertices in the same layer, then G has a cycle; otherwise, G has no cycles.

Prove that this algorithm is correct, or give a counterexample.

Question 4 (20 points). Recall *Dijkstra's Algorithm* for finding shortest paths in a directed, weighted graph.

- (a) Why doesn't Dijkstra's algorithm work if edges in the graph can have negative weights? Give an example of a graph where Dijkstra's algorithm fails to find the minimum-weight path between a pair of vertices.
- (b) What happens if we replace the word "smallest" in Dijkstra's algorithm with the word "biggest"—that is, we use a max priority queue so we pull out the *biggest* edge on each iteration, and when we add a vertex u to the set S , for each edge (u, v) we update $d[v]$ to be the *maximum* of $d[v]$ and the sum $d[u] + w_{uv}$. Can we use this modified Dijkstra's algorithm to find *longest* paths between nodes in a graph? Prove that this works, or give a counterexample where it doesn't.

Question 5 (30 points). Let $G = (V, E)$ be an undirected, connected graph with non-negative edge lengths. We say that T is a *min-max spanning tree* on G if T is a spanning tree whose longest edge is as short as possible. That is, if we imagine taking all the possible spanning trees of G and sorting them by the length of their *longest* edge, in increasing order, T would be first (or tied for first) on the list. Put yet another way, T is a min-max spanning tree if and only if every other spanning tree T' has some edge which is longer than (or equal to) every edge in T .

Prove or disprove the following claims:

- (a) If T is a minimum spanning tree for G , then T is a min-max spanning tree for G .
- (b) If T is a min-max spanning tree for G , then T is a minimum spanning tree for G .