

## Algorithms: Applications of BFS

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Suppose we have a graph  $G = (V, E)$ . A given graph could have few edges, or lots of edges, or anything in between. Let's think about the range of possible relationships between  $V$  and  $E$ .

1 How big can  $|E|$  be, relative to  $|V|$ ?

(a) The smallest possible value of  $|E|$  is \_\_\_\_\_.

(b)  $|E|$  is  $O(\quad)$  because \_\_\_\_\_.

(c) When  $G$  is a tree,  $|E|$  is  $\Theta(\quad)$  because \_\_\_\_\_.

Now, recall from last class that we showed breadth-first search (BFS) can be implemented to run in  $\Theta(|E|)$  time.

2 In terms of  $\Theta$ , how fast does BFS run, as a function of  $|V|$ , when  $G$  is a tree?

3 How fast does BFS run, as a function of  $|V|$ , when  $G$  is very dense, *i.e.* it contains some constant fraction (say, half) of all possible edges?

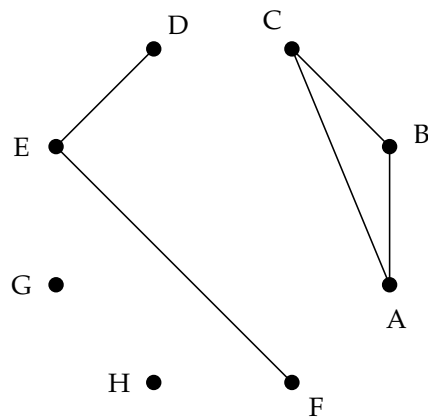
### A first application of BFS

- 4 Describe an algorithm to find the connected components of a graph  $G$ .

**Input:** a graph  $G = (V, E)$

**Output:** a set of sets of vertices,  $\text{Set}\langle\text{Set}\langle\text{Vertex}\rangle\rangle$ , where each set contains the vertices in some (maximal) connected component. That is, all the vertices within each set should be connected; no vertex should be connected to vertices in any other sets; and every vertex in  $V$  should be contained in exactly one of the sets.

For example, given the graph below, the algorithm should return  $\{\{D, E, F\}, \{C, B, A\}, \{G\}, \{H\}\}$ .



Describe your algorithm (using informal prose or pseudocode) and analyze its asymptotic running time.



## *A second application of BFS*

### *Model 1: Directed graphs*

See the board for examples of *directed* graphs.

- 5 What is the difference between directed graphs and the (undirected) graphs we saw on a previous activity?
- 6 The previous activity defined graphs as consisting of a set  $V$  of vertices and a set  $E$  of edges, where each edge is a set of two vertices. How would you modify this definition to allow for directed graphs?
- 7 For each of the following graph terms/concepts, say whether you think its definition needs to be modified for directed graphs; if so, say what the new definition should be.
  - 1 *vertex*
  - 2 *degree*
  - 3 *path*
  - 4 *cycle*



- 8 What (if anything) about our implementation of BFS needs to be modified for BFS to work sensibly on directed graphs?

**Definition 1.** A directed graph  $G = (V, E)$  is *strongly connected* if for any two vertices  $u, v \in V$  there is a (directed) path from  $u$  to  $v$ , and also from  $v$  to  $u$ .

- 9 Describe a brute force algorithm for determining whether a given directed graph  $G$  is strongly connected.
- 10 Analyze the running time of your algorithm in terms of  $\Theta$ .



Let's see if we can do better!

**Theorem 2.** *A directed graph  $G = (V, E)$  is strongly connected if and only if for any vertex  $s \in V$ , every other vertex in  $G$  is mutually reachable with  $s$  (that is, for each  $v \in V$  there is a directed path from  $s$  to  $v$  and another directed path from  $v$  to  $s$ ).*

11 In order to prove this “if and only if” statement, we must prove

both \_\_\_\_\_

and \_\_\_\_\_.

*Proof.*

Hint: draw a picture!

□



**Definition 3.** Given a directed graph  $G$ , its *reverse graph*  $G^{\text{rev}}$  is the graph with the same vertices and all edges reversed.

**Theorem 4.** A directed graph  $G = (V, E)$  is strongly connected if and only if given any  $s \in V$ , all vertices are reachable from  $s$  in  $G$ , and all vertices are reachable from  $s$  in  $G^{\text{rev}}$ .

*Proof.*

Hint: what is the relationship between this theorem and Theorem 2?

□

- 12 Based on the above theorem, describe an algorithm to determine whether a given directed graph  $G = (V, E)$  is strongly connected, and analyze its running time.

