Using lhs2TeX and sproof

January 12, 2016

This document explains how to make use of

- lhs2TeX, for typesetting Haskell code in LATEX, and
- the sproof environment, for typesetting structured proofs.

The document itself explains most of what you need to know, but for even more insight you can take a look at the IAT_EX source used to create this document as well.

1 Typesetting Haskell with lhs2TeX

To install lhs2TeX, it should suffice to issue the command

cabal install lhs2TeX

To use it, start by creating a IATEX document within a .lhs file (if you use emacs, you can tell it to use latex-mode instead of literate Haskell mode by putting % -*- mode: LaTeX -*- at the top of your file). At the top of your file, immediately following the \documentclass command, you should put

%include polycode.fmt

which will instruct lhs2TeX to import a standard library of tools and definitions.

Within your LATEX document, you can include snippets of Haskell code by enclosing them in vertical bars. For example,

llength . map f == length|

gets typeset as $length \circ map \ f \equiv length$. Note that there are several styles for typeset Haskell code. The default is to use italics and nice "mathy" symbols like \circ in place of ., \equiv in place of ==, \leq in place of <=, and so on. This looks nice but can sometimes be harder to read if you are not used to it. If you would rather have good old typewriter font used for your Haskell code, you can pass the --verb option to lhs2TeX.

You can also include larger sections of Haskell code by enclosing them in

\begin{code} ... \end{code}

 or

\begin{spec} ... \end{spec}

(The only difference between the two is that anything in a spec environment will be ignored when the file is loaded into ghci or compiled with ghc.) For example,

```
\begin{code}
map :: (a -> b) -> [a] -> [b]
map f [] = []
map f (x:xs) = f x : map f xs
\end{code}
```

is typeset as

```
 \begin{array}{l} map :: (a \rightarrow b) \rightarrow [a] \rightarrow [b] \\ map \ f \ [] \qquad = [] \\ map \ f \ (x : xs) = f \ x : map \ f \ xs \end{array}
```

To achieve the nice vertical alignment of the equals signs (or anything else), vertically align them in the source code, taking care to leave at least *two* spaces before each of the things to be aligned (note how map f $(x:xs) = f x \dots$ actually has two spaces before the = sign).

 $\tt lhs2TeX$ acts as a preprocessor that takes . <code>lhs</code> files and produces <code>.tex</code> files. You can use it with a command line like

```
lhs2TeX myfile.lhs > myfile.tex
```

followed by compiling the .tex file normally (*e.g.* with pdflatex). Recall that you can also pass the --verb option to lhs2TeX to use a typewriter font for your Haskell code.

There is much, much more you can do with lhs2TeX, including automatically typesetting variable names like x2 as x_2 , introducing your own custom formatting for certain functions or operators, and automatically evaluating Haskell expressions via ghci and inserting their output into your document. For more information, see http://www.andres-loeh.de/lhs2tex/. It appears there is as of yet no manual that specifically accompanies the latest release (1.19), but the 1.17 manual should work just fine (http://www.andres-loeh.de/lhs2tex/Guide2-1.17.pdf).

2 Typesetting structured proofs with sproof

The first step is to download sproof.sty. You can put this file somewhere IAT_EX knows to look for it, or (much simpler) just put a copy of it in the same folder as any .tex files you want to use it with. To use it, include \usepackage{sproof} in your document preamble.

Structured proofs are created with the sproof environment. Each expression or statement is typeset with the \stmt{...} command, which automatically places its contents in math mode. In between the expressions/statements, you can use the \reason{...}{...} command. The first argument to \reason is a transitive binary operator (in math mode); the second argument is a justification (in text mode). So, for example,

```
\begin{sproof}
  \stmt{1 + 1 + 1}
  \reason{=}{arithmetic}
  \stmt{1 + 2}
  \reason{=}{more arithmetic}
  \stmt{3}
  \reason{\leq}{duh}
  \stmt{5}
 \end{sproof}
```

produces

 $\begin{array}{c} 1+1+1 \\ = \\ 1+2 \\ = \\ 3 \\ \leqslant \\ 5 \end{array} \qquad \left\{ \begin{array}{c} \text{arithmetic} \end{array} \right\} \\ \left\{ \begin{array}{c} \text{more arithmetic} \end{array} \right\} \\ \left\{ \begin{array}{c} \text{duh} \end{array} \right\} \end{array}$

As another example, consider typesetting an inductive proof of the fact that for all lists xs and functions f, length $(map \ f \ xs) = length \ xs$. The empty list case is simple enough; consider the case when xs is a cons, y : ys. As our inductive hypothesis, we get to assume

$$length (map f ys) = length ys,$$

and we must show the same holds for y : ys. We reason as follows:

length (map f (y:ys)) $\begin{cases} length (f y : map f ys) \\ 1 + length (map f ys) \\ \end{cases}$ { defn of map= } defn of *length* } = IH } $1 + length \ ys$ { defn of *length* } _ length (y:ys)